

ResilientDB: Global Scale Resilient Blockchain Fabric



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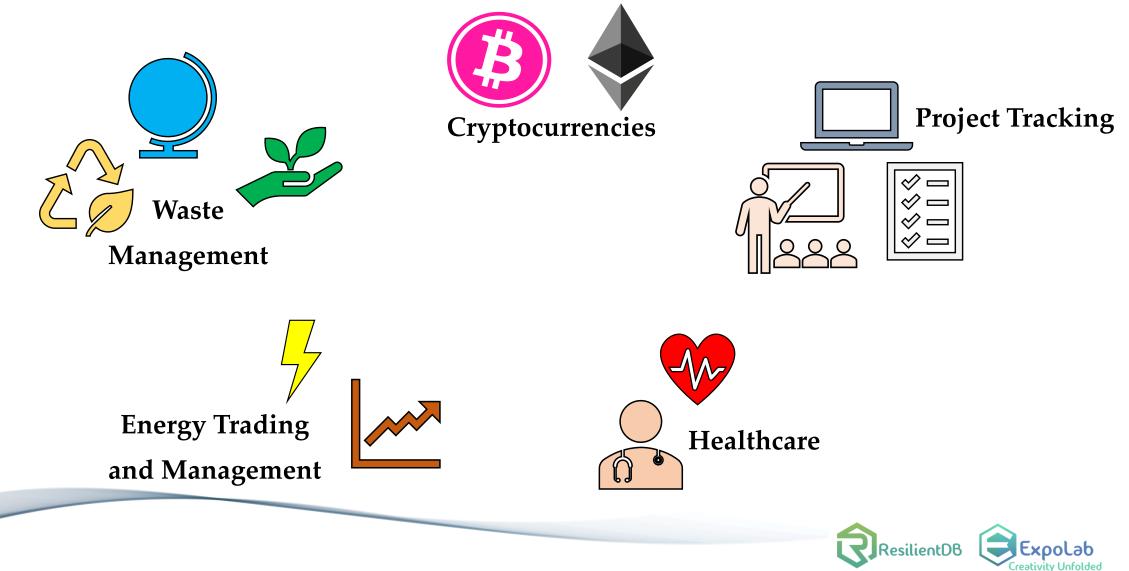
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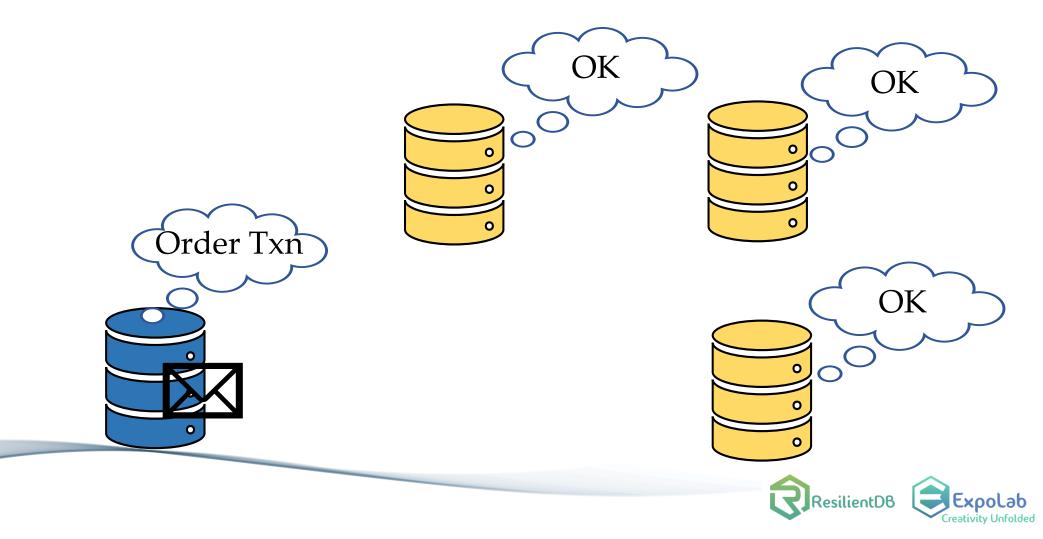




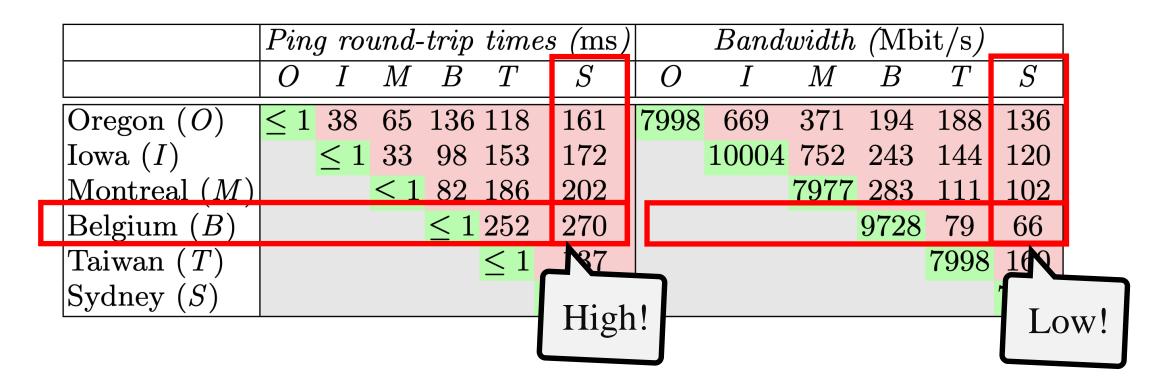
Blockchain Applications



At the core of *any* Blockchain application is a Byzantine Fault-Tolerant (BFT) consensus protocol.



Challenges For Geo-Scale Blockchains



Blockchain expects decentralization \rightarrow Replicas maybe geo-distributed



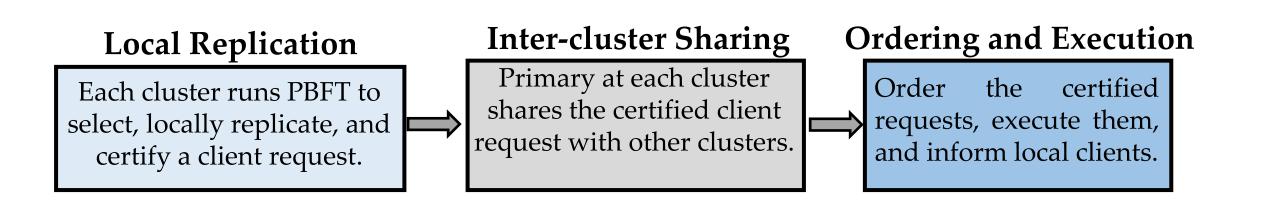
Limitations of Existing Consensus Protocols

Protocol	Decisions	Communication		Centralized
		(Local)	(Global)	
GEOBFT (our paper)	Z	$\mathcal{O}(2\mathbf{zn}^2)$	$\mathcal{O}(\mathbf{fz}^2)$	No
\lor single decision	1	$\mathcal{O}(4\mathbf{n}^2)$	$\mathcal{O}(\mathbf{fz})$	No
STEWARD	1	$\mathcal{O}(2\mathbf{zn}^2)$	$\mathcal{O}(\mathbf{z}^2)$	Yes
Zyzzyva	1	$\mathcal{O}(\mathbf{zn})$		Yes
PBFT	1	$\mathcal{O}(2(\mathbf{zn})^2)$		Yes
PoE	1	$\mathcal{O}((\mathbf{zn})^2)$		Yes
HotStuff	1	$\mathcal{O}(8(\mathbf{zn}))$		Partly

Normal-case metrics for a system with z clusters, each with n replicas of which at most f , n > 3f , are Byzantine



GeoBFT Protocol

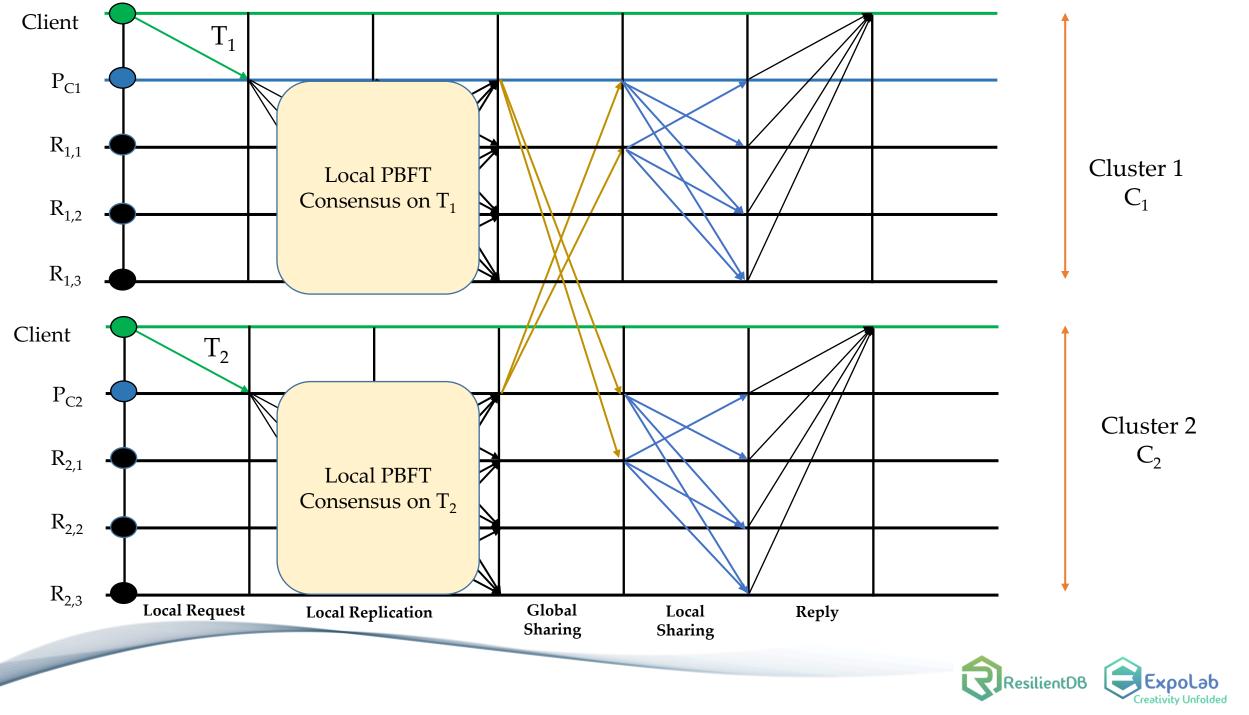


Topology-aware protocol.

Parallel consensus at each cluster.

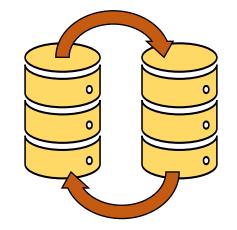
Low inter-cluster communication overheads.





Local Replication

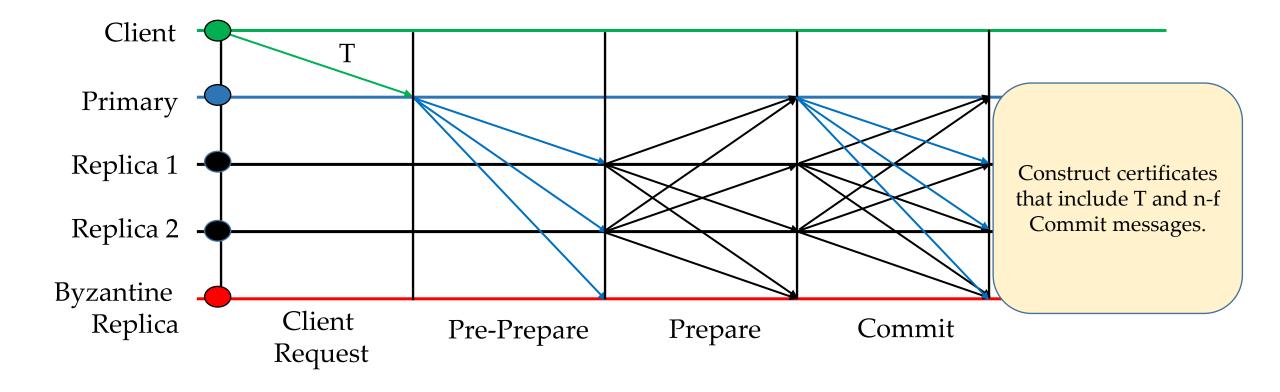
- GeoBFT employs PBFT for local replication.
- Tolerates up to **f** failure out of **3f+1** replicas



- Three phases of which two require quadratic communication complexity.
- Safety is always guaranteed and Liveness is guaranteed in periods of partial synchrony.
- View-Change protocol for replacing malicious primary

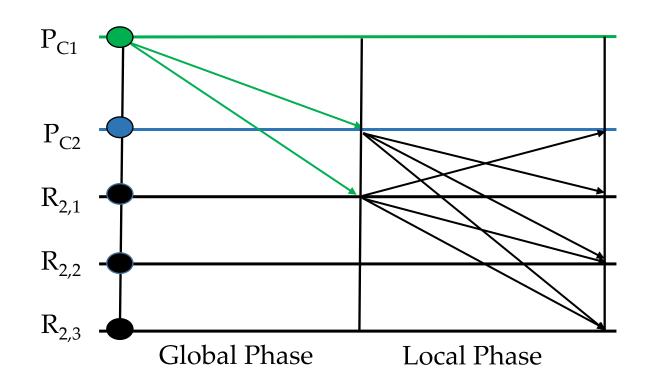


PBFT Civil Execution





Inter-Cluster Sharing



The Primary P_{C1} sends a certificate that includes the client request and commit

messages from n-f replicas of Cluster C_1 .



Ordering and Execution

GeoBFT orders requests deterministically.

For i < j, requests of Cluster C_i are executed before requests of cluster C_i .

For example: requests of C_1 are executed before C_2 .



Implementation on ResilientDB



ResilientDB associates a multi-threaded deep-pipelined architecture with each replica.



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ResilientDB is open-sourced at <u>https://resilientdb.com/</u>

Ledger (Blockchain) Management

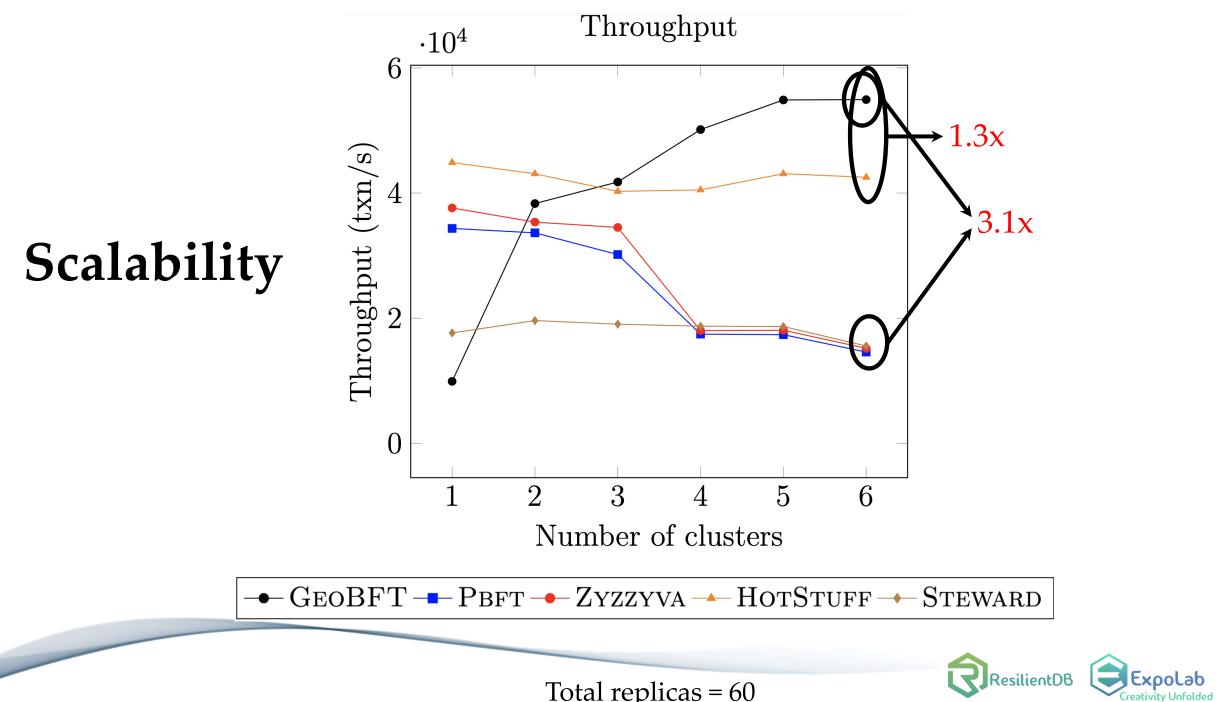
- ith block in the ledger contains the ith executed request.
- In each round of GeoBFT, each replica executes z requests, each belonging to a different cluster C_i, 1 <= i <= z.
- Hence, in each round, each replica creates z blocks.
- To ensure immutability, each block includes both client requests and exchanged certificates.



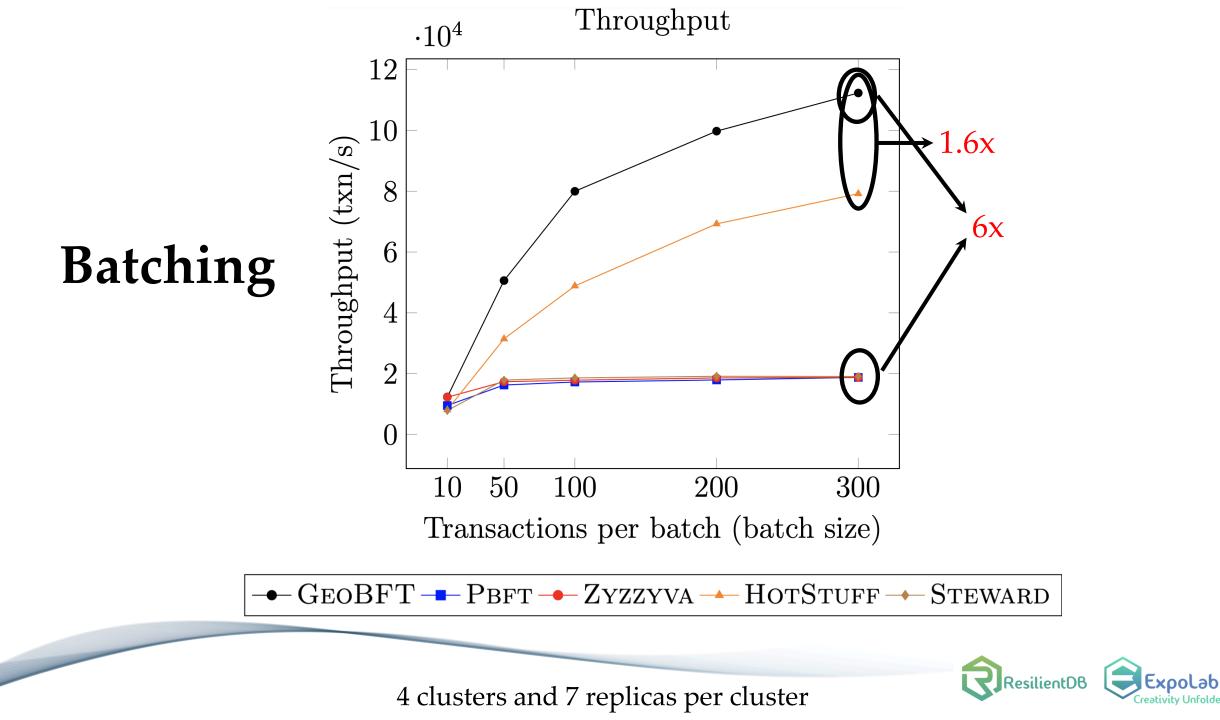
Evaluation on ResilientDB

- Google cloud used for deploying replicas and clients.
- Each replica used 8-core Intel Skylake CPUs and had access to 16 GB memory.
- Total 160K clients deployed on eight 4-core machines.
- Workload provided by Yahoo Cloud Serving Benchmark (YCSB).
- Replicas deployed across six different regions: Oregon, Iowa, Montreal, Belgium, Taiwan and Sydney.
- Primaries for centralized protocol placed at Oregon (highest bandwidth).





Total replicas = 60



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Conclusions and Final Remarks

- For achieving faster local replication, GeoBFT can employ other efficient BFT protocols, such as PoE.
- Modern cryptographic techniques such as Threshold signatures can be used in place of sending n-f Commit messages.
- If a cluster does not have a request, it can send "no-op" messages.
- GeoBFT optimizes consensus by reducing global communication costs.
- Parallel local replication helps to increase system throughput.



References

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